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Problems of cryptography as discrete optimization tasks

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Abstract

In this contribution problems encountered in the field of cryptology, are introduced as discrete optimization tasks. Two evolutionary computation algorithms, namely the particle swarm optimization method and the differential evolution method, are applied to handle these problems. The results indicate that the dynamic of this type of discrete optimization problems makes it difficult for the methods to retain information.

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1. Introduction

The field of cryptography has motivated a number of hard and complex computational problems. Such problems are the integer factorization problem related to the RSA cryptosystem; the index computation or the discrete logarithm problem related to the El Gamal cryptosystem, as well as, to the Diffie–Hellman key exchange and others [2,3,9,10]. The assumption that these problems are in general computationally intractable in polynomial time forms the basis of the reliability of most contemporary cryptosystems.

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In this paper, a number of problems originating from the integer factorization problem are formulated as discrete optimization tasks. The integer factorization problem forms the basis of the RSA cryptosystem. Thus, optimization methods that can efficiently tackle the problems under consideration in acceptable computational time, could be regarded as a useful tool for cryptanalysis of the corresponding type of cryptosystems.

Evolutionary computation algorithms are stochastic optimization methods inspired by either natural evolution or social behavior. Genetic algorithms (GA) [4], evolution strategies (ES) [11], the differential evolution algorithm (DE) [12], as well as, the particle swarm optimization (PSO) [1,5] belong to this class of methods. All the aforementioned optimization algorithms are designed to address problems involving discontinuous and multimodal objective functions, the existence of numerous local minima, constrained optimizations tasks, and disjoint search spaces [4,5,11]. Optimization techniques for real search spaces can be applied to discrete optimization problems with minor modifications. A straightforward approach is to round off the optimum solution to the nearest integer [6,8].

Two evolutionary computation algorithms, namely the particle swarm optimization method and the differential evolution algorithm, are applied to tackle several instances of the proposed optimization problems. The performance of both these methods is compared with simple random search.

The rest of this contribution is organized as follows. The definition of the problems along with the transformations to discrete optimization tasks, are given in Section 2. In Section 3, experimental setup of the considered methods and results are reported. Conclusions are given in Section 4.

2. Problem formulation

The first problem under consideration is defined as follows: given a composite integer N, find pairs of x, $y \in \mathbb{Z}_N^*$, such that $x^2 \equiv y^2 \pmod{N}$, with $x \not\equiv \pm y \pmod{N}$. This problem is equivalent to finding non-trivial factors of N, as N divides $x^2 - y^2 = (x - y)(x + y)$, but N does not divide either x - y or x + y. Hence the gcd(x - y, N) is a non-trivial factor of N (random square factorization algorithm) [7].

The prescribed problem can be formulated as a discrete optimization task by defining the minimization function $f : \{1, ..., N - 1\} \times \{1, ..., N - 1\} \mapsto \{0, ..., N - 1\}$, with

$$f(x, y) = x^2 - y^2 \pmod{N},$$

subject to the constraints $x \neq \pm y \pmod{N}$. The constraint x = -y can be incorporated to the problem by changing the domain of the function. Thus, the problem reduces to minimizing the function $g: \{2, 3, ..., (N-1)/2\} \times \{2, 3, ..., (N-1)/2\} \mapsto \{0, ..., (N-1)\}$, with

$$g(x, y) = x^2 - y^2 \pmod{N},$$

subject to the constraint $x \neq y \pmod{N}$. The minimization problem is two-dimensional and the global minimum of the function g is zero. A plot of the function f(x, y) for N=5*7=35 is given in Fig. 1(a) and the contour plot for the value f(x, y) = 0 is given in Fig. 1(b).

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Fig. 1. (a) Plot of function $f(x, y) = x^2 - y^2 \pmod{N}$, for N = 35 (left), and (b) contour plot of function $f(x, y) = x^2 - y^2 \pmod{N}$, for N = 35 at value 0 (right).

Similar problems that can be studied are the following: Minimize the function $h: \{1, ..., N-1\} \mapsto \{0, ..., N-1\}$, with

 $h(x) = (x - a)(x - b) \pmod{N},$

where *a*, *b* non-zero integers and $x \neq a \pmod{N}$, $x \neq b \pmod{N}$. As an example of this problem, we have considered the minimization of the function $h_e(x) = (x - 1)(x - 2)$, where $x \neq 1 \pmod{N}$ and $x \neq 2 \pmod{N}$. In a more general form one can consider the minimization of the function

$$w(x) = (x - a)(x - b) \cdots (x - m) \pmod{N},$$

where $x \in \{0, ..., N-1\}$ and $x \not\equiv \{a, b, ..., m\} \pmod{N}$. As an example of this problem, we have studied the function $w_e(x) = (x+1)(x-1)(x-2) \pmod{N}$, with $x \not\equiv \{-1, 1, 2\} \pmod{N}$.

3. Experimental setup and results

The PSO [1] and DE [12] methods, were applied on the considered problems along with the random search technique. The global and local PSO variants of both the inertia weight and the constriction factor versions, as well as the *DE/rand/1/bin* and *DE/best/2/bin* variants of the DE algorithm, have been used. For both the PSO variants typical parameter values were used, while the size of the neighborhood for the local variant of the PSO was taken equal to 1. Preliminary experiments indicated that the value of maximum velocity V_{max} of the PSO's particles affects its performance significantly. The value $V_{max} = \lfloor (UpBound - LoBound)/5 \rfloor$, where UpBound denotes the upper bound of the function's domain and LoBound the lower bound of the function's domain, produced the most promising results and therefore it was adopted in all the experiments. For the DE algorithm, the parameters were set at the values F = 0.5 and CR = 0.5. All populations were constrained in the feasible region of the corresponding problem.

Table 1Results for the minimization of function g

| Ν | Method | Suc. rate (%) | Mean FE | SD FE | Median FE | Min FE |
|----------------|------------|---------------|------------|------------|------------|--------|
| N = 199 * 211 | PSOGW | 56 | 8844.643 | 5992.515 | 8325.000 | 660 |
| | PSOGC | 48 | 7149.375 | 5272.590 | 5355.000 | 330 |
| | PSOLW | 51 | 8329.412 | 6223.142 | 7050.000 | 270 |
| | PSOLC | 51 | 7160.588 | 6001.276 | 5940.000 | 420 |
| | DE1 | 4 | 517.500 | 115.866 | 465.000 | 450 |
| | DE2 | 9 | 5476.667 | 6455.651 | 1830.000 | 60 |
| | RS | 66 | 9104.015 | 5862.358 | 8700.500 | 22 |
| N = 293 * 307 | PSOGW | 41 | 16210.244 | 11193.375 | 15090.000 | 120 |
| | PSOGC | 45 | 16818.667 | 12664.632 | 13800.000 | 630 |
| | PSOLW | 58 | 18455.690 | 12870.897 | 14520.000 | 270 |
| | PSOLC | 50 | 16374.000 | 13597.782 | 13365.000 | 120 |
| | DE1 | 7 | 1598.571 | 1115.488 | 1470.000 | 120 |
| | DE2 | 19 | 17815.263 | 12484.580 | 16290.000 | 2730 |
| | RS | 64 | 21548.531 | 13926.751 | 20852.500 | 57 |
| N = 397 * 401 | PSOGW | 53 | 31965.849 | 24423.975 | 27570.000 | 780 |
| | PSOGC | 45 | 32532.667 | 22652.983 | 33210.000 | 1740 |
| | PSOLW | 55 | 31472.182 | 23394 791 | 22620.000 | 720 |
| | PSOLC | 54 | 38156 111 | 22925 970 | 37665.000 | 750 |
| | DE1 | 1 | 1680,000 | 0.000 | 1680.000 | 1680 |
| | DE1 DE2 | 12 | 27722 500 | 17498 736 | 28620.000 | 180 |
| | RS | 60 | 27302.567 | 21307.031 | 23607.500 | 145 |
| N = 499 * 503 | PSOGW | 56 | 49893 750 | 37515 327 | 44640 000 | 930 |
| 11 = 199 + 505 | PSOGC | 55 | 49975 636 | 36727 380 | 41760.000 | 300 |
| | PSOLW | 55 | 49207 091 | 34053 904 | 50430.000 | 2010 |
| | PSOLC | 46 | 48443 478 | 34677 039 | 43470.000 | 1920 |
| | DE1 | 1 | 2480.000 | 0.000 | 2480.000 | 2480 |
| | DE2 | 8 | 67245.000 | 35114 316 | 64770.000 | 14730 |
| | RS | 61 | 54139.443 | 38642.970 | 48743.000 | 140 |
| N = 599 * 601 | PSOGW | 52 | 72175.000 | 48653.823 | 71550.000 | 600 |
| | PSOGC | 51 | 81476.471 | 53666.543 | 75100.000 | 5000 |
| | PSOLW | 49 | 78651.020 | 48197.105 | 67400.000 | 11200 |
| | PSOLC | 52 | 69542.308 | 48837.949 | 53050.000 | 2500 |
| | DE1 | 2 | 4700.000 | 4808 326 | 4700.000 | 1300 |
| | DE2 | 5 | 8620.000 | 8078.180 | 9300.000 | 800 |
| | RS | 64 | 86123.656 | 47504.284 | 89392.500 | 904 |
| N = 691 * 701 | PSOGW | 46 | 207443.478 | 163585.340 | 214800.000 | 800 |
| | PSOGC | 46 | 175426.086 | 138118.794 | 149200.000 | 800 |
| | PSOLW | 60 | 196993.334 | 146204.518 | 144500.000 | 9200 |

| Ν | Method | Suc. rate (%) | Mean FE | SD FE | Median FE | Min FE |
|---|--------|---------------|------------|------------|------------|--------|
| | PSOLC | 52 | 209307.692 | 163833.606 | 200100.000 | 1800 |
| | DE1 | 2 | 23800.000 | 25000.000 | 23800.000 | 21000 |
| | DE2 | 10 | 71000.000 | 95357.642 | 15200.000 | 1600 |
| | RS | 60 | 185932.334 | 126355.926 | 154999.000 | 2828 |

Table 1 (continued)

Table 2 Results for the minimization of the functions h_e and w_e , for N = 103 * 107

| Function | Method | Suc. rate (%) | Mean FE | SD FE | Median FE | Min FE |
|----------------|--------|---------------|----------|----------|-----------|--------|
| he | PSOGW | 51 | 2013.333 | 1483.535 | 1500.000 | 100 |
| | PSOGC | 57 | 1974.035 | 1609.228 | 1420.000 | 60 |
| | PSOLW | 59 | 1677.288 | 1254.688 | 1420.000 | 60 |
| | PSOLC | 58 | 2385.862 | 1676.898 | 2040.000 | 120 |
| | DE1 | 1 | 100.000 | 0.000 | 100.000 | 100 |
| | DE2 | 1 | 80.000 | 0.000 | 80.000 | 80 |
| | RS | 65 | 2099.646 | 1448.007 | 2056.000 | 6 |
| w _e | PSOGW | 79 | 1382.785 | 1265.927 | 820.000 | 40 |
| | PSOGC | 84 | 1402.857 | 1442.194 | 930.000 | 40 |
| | PSOLW | 80 | 1757.750 | 1544.267 | 1110.000 | 40 |
| | PSOLC | 85 | 1416.000 | 1329.034 | 880.000 | 40 |
| | DE1 | 1 | 60.000 | 0.000 | 60.000 | 60 |
| | DE2 | 1 | 80.000 | 0.000 | 80.000 | 80 |
| | RS | 96 | 1507.969 | 1328.913 | 1104.000 | 7 |

For the minimization of the function g, the performance of the methods was investigated for several instances of N, from the value N = 199 * 211 = 41,989 up to N = 691 *701 = 484,391. For each N considered, 100 independent runs were performed and the corresponding results are exhibited in Table 1. Concerning the notation used in the Table, PSOGW corresponds to the global variant of PSO method with inertia weight; PSOGC is the global variant of PSO with constriction factor; PSOLW is PSO's local variant with inertia weight; PSOLC is PSO's local variant with constriction factor, DE1 corresponds to the *DE/rand/1/bin* and DE2 to the *DE/best/2/bin* variants of DE method. Random search results are denoted as RS. A run is considered to be successful if the algorithm identifies the global minimizer within a prespecified number of function evaluations. The function evaluations threshold was taken equal to the cardinal of integers in the domain of the function studied. The success rates of each algorithm, that is the proportion of the times it achieved the global minimizer within the prespecified threshold, the minimum number, the median, the mean value and the standard deviation of function evaluations (FE) needed for success, are reported. The results indicate that the variants of PSO method outperform the variants of the DE method over these instances and with this parameter setup. Moreover, the performance of the DE method decreases as the value of *N* increases in contrast to PSO which appears to be more stable with respect to this parameter. However, in contrast to the known behavior of the evolutionary computation methods, the random search technique outperforms both these methods and their variants. This fact suggests that the almost random behavior of the specific kind of problems makes it quite difficult for the methods to retain knowledge about their dynamics. Similar results were reported on the minimization of the functions h_e and w_e . For N = 103 * 107 the results are reported in Table 2.

4. Conclusions

In this paper, a number of problems originating from the integer factorization problem are introduced as discrete optimization tasks. Since the integer factorization problem forms the basis of many contemporary cryptosystems, optimization methods that can efficiently and effectively tackle the considered problems, could constitute a useful tool for the cryptanalysis of the corresponding cryptosystems.

Two evolutionary computation algorithms, namely the particle swarm optimization method and the differential evolution method, are applied to handle these problems. Their results, compared with the random search technique, indicate that this special kind of discrete optimization problems seem to have a dynamic that makes it difficult for the methods to retain information.

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